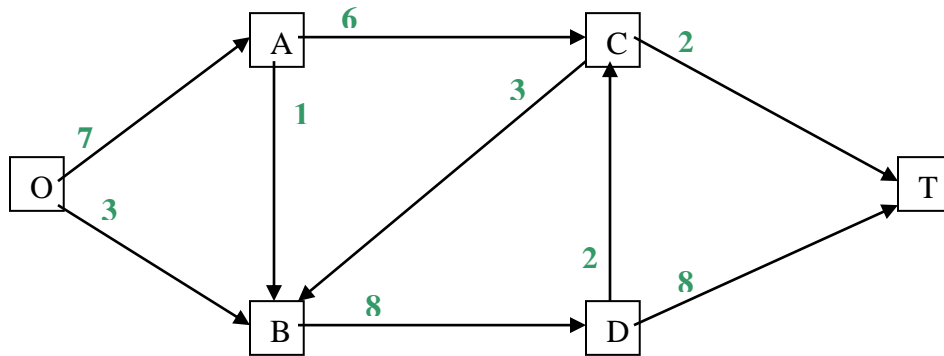


Solutions

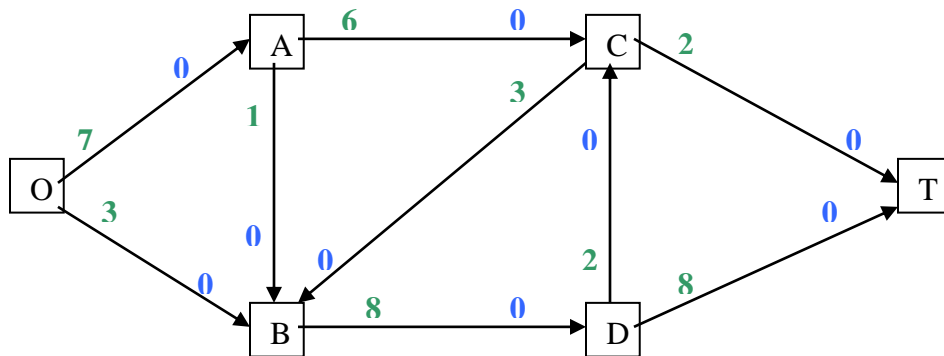
Maximum Flow Problem.

You are given the following directed network with source O and sink T.

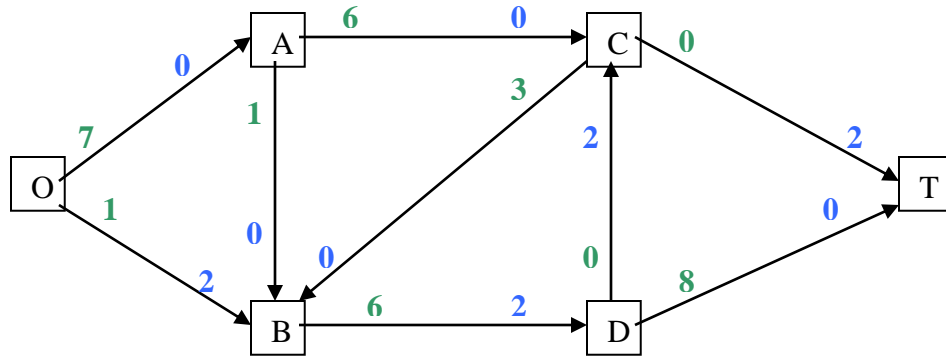


- a) Find a maximum flow from O to T in the network (choose $O \rightarrow B \rightarrow D \rightarrow C \rightarrow T$ as the first augmenting path).

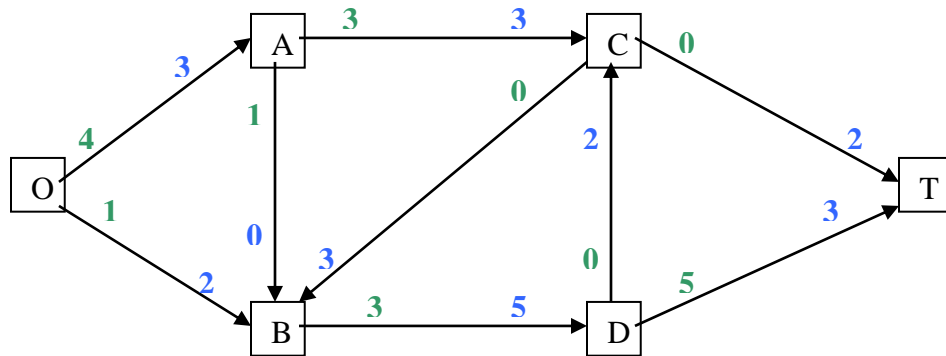
The first residual network:



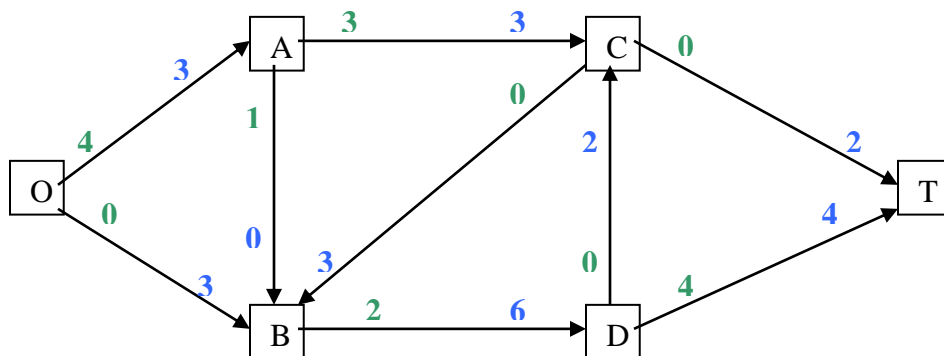
Iteration 1: As required, $O \rightarrow B \rightarrow D \rightarrow C \rightarrow T$ is chosen as the first augmenting path. (The reason I asked you to start from this path is that it creates necessity for reversing some flow in later iterations.) Its residual capacity is $\min\{3, 8, 2, 2\} = 2$. After sending 2 units of flow through the path, the resulting residual network is:



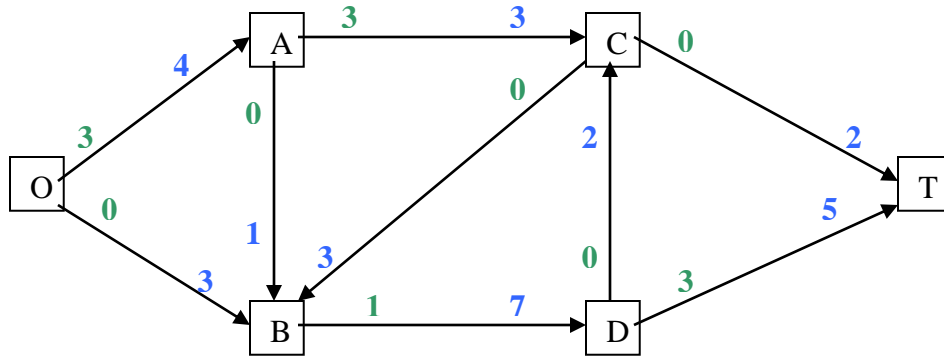
Iteration 2: Choose $O \rightarrow A \rightarrow C \rightarrow B \rightarrow D \rightarrow T$ as the next augmenting path. Its residual capacity is $\min\{7,6,3,6,8\}=3$. After sending 3 units of flow through the path, the new residual network is:



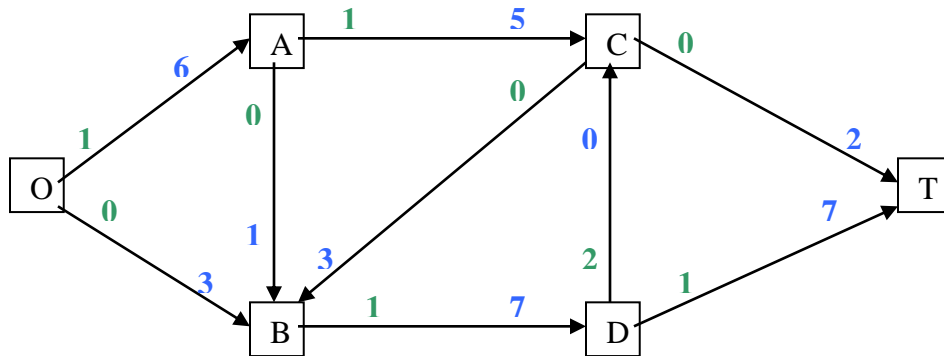
Iteration 3: Choose $O \rightarrow B \rightarrow D \rightarrow T$ as the next augmenting path. Its residual capacity is $\min\{1,3,5\}=1$. After sending 1 units of flow through the path, the new residual network is:



Iteration 4: Choose $O \rightarrow A \rightarrow B \rightarrow D \rightarrow T$ as the next augmenting path. Its residual capacity is $\min\{4,1,2,4\}=1$. After sending 1 units of flow through the path, the new residual network is:



Iteration 5: Choose $O \rightarrow A \rightarrow C \rightarrow D \rightarrow T$ as the next augmenting path. Its residual capacity is $\min\{3, 3, 2, 3\} = 2$. (Note that we reverse the flow on arc $D \rightarrow C$). After sending 2 units of flow through the path, the new residual network is:



No more augmenting paths from O to T are left. Thus, the current residual network is optimal. The blue numbers on the arcs show the optimal flow values. Maximum flow value is $7 + 2 = 9$ (the flow amount entering the sink).

b) Find a minimum cut. What is its capacity?

The nodes that can be reached from the source by augmenting paths are A and C. Thus, the O-side of the minimum cut is $\{O, A, C\}$. The minimum cut is

$$\text{MinCut} = \{O \rightarrow B, A \rightarrow B, C \rightarrow B, C \rightarrow T\}$$

(Note that arc $D \rightarrow C$ is not in the cut because it goes from T-side to O-side).

The capacity of the minimum cut is $3 + 1 + 3 + 2 = 9$ which is equal to the maximum flow value.

The minimum cut is shown in the network below.

