

## Unbounded Problems

$$\begin{aligned}
 \max \quad & x_1 + x_2 \\
 \text{s.t.} \quad & -x_1 + x_2 \leq 1 \\
 & x_1 - x_2 \leq 1 \\
 & x_1, x_2 \geq 0
 \end{aligned}$$

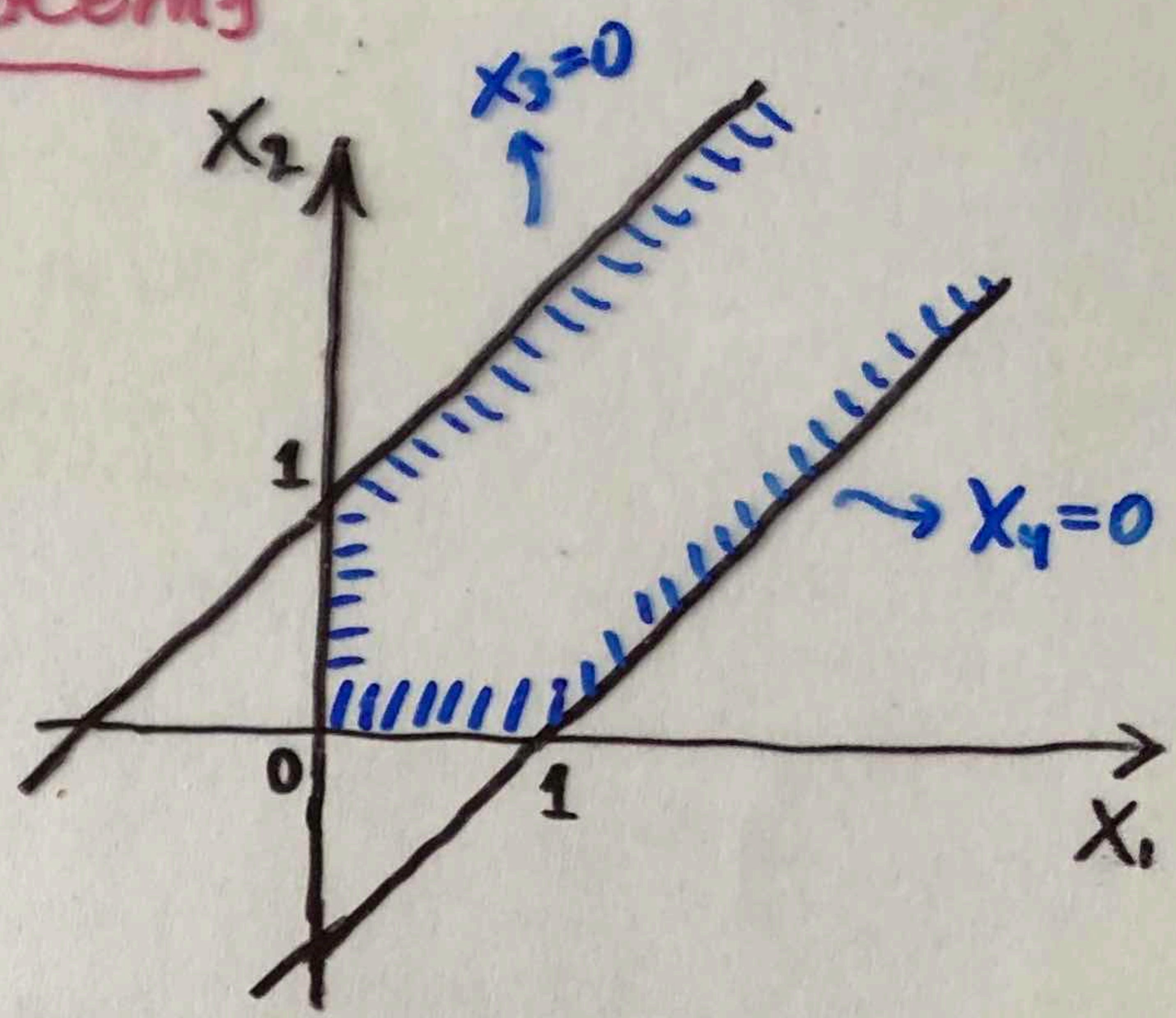


Tableau (after introducing slack variables  $x_3$  and  $x_4$ )

Basic var.	$x_1$	$x_2$	$x_3$	$x_4$	RHS
$z$	-1	-1	0	0	0
$x_3$	-1	1	1	0	1
$x_4$	1*	-1	0	1	1

$B = \{3, 4\}, N = \{1, 2\}$

Basic var.	$x_1$	$x_2$	$x_3$	$x_4$	RHS
$z$	0	-2	0	1	1
$x_3$	0	0	1	1	2
$x_1$	1	-1	0	1	1

$B = \{3, 1\}, N = \{4, 2\}$

$x_2$  is the only candidate to enter the basis; but we can increase  $x_2$  infinitely without violating the constraints.

Thus, the problem is unbounded.

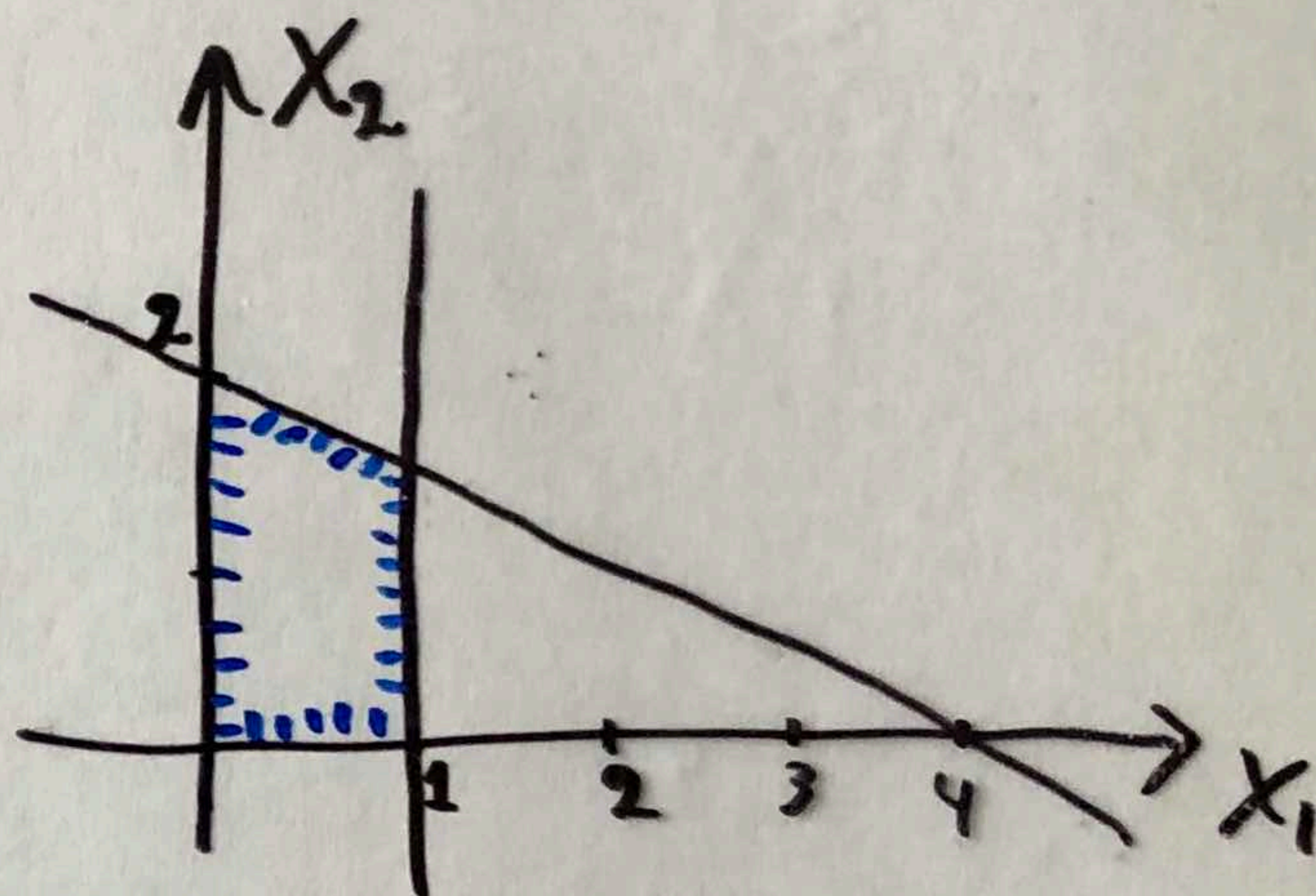
(2)

$$x(\lambda) = \begin{pmatrix} 1 \\ 0 \\ 2 \\ 0 \end{pmatrix} + \lambda \cdot \begin{pmatrix} 1 \\ 1 \\ 0 \\ 0 \end{pmatrix} \text{ is a feasible ray} \\ \text{(feasible for all } \lambda \geq 0 \text{);}$$

and  $Z = 1 + 2\lambda$  on that ray.

### Multiple Optimal Solutions

$$\begin{aligned} \max \quad & X_1 + 2X_2 \\ \text{s.t.} \quad & X_1 + 2X_2 \leq 4 \\ & X_1 \leq 1 \\ & X_1, X_2 \geq 0 \end{aligned}$$



After introducing slack variables:

Basic var.	$X_1$	$X_2$	$X_3$	$X_4$	RHS
$Z$	-1	-2	0	0	0
$X_3$	1	2*	1	0	4
$X_4$	1	0	0	1	1

Basic var.	$X_1$	$X_2$	$X_3$	$X_4$	RHS
$Z$	0	0	1	0	4
$X_2$	$\frac{1}{2}$	1	$\frac{1}{2}$	0	2
$X_4$	1	0	0	1	1

$$B = \{2, 4\}, N = \{1, 3\}$$

$$X^* = \begin{pmatrix} 0 \\ 2 \\ 0 \\ 1 \end{pmatrix} \text{ is opt.}$$

$$Z^* = 4$$

- In this tableau,  $x_1$  is a nonbasic variable and its coefficient is 0 in row 0  $\rightarrow$  increasing  $x_1$  will not change the obj. f-n value.
- Do another iteration, entering  $x_1$  into basis. Min-ratio test:  $\min(\frac{2}{\frac{1}{2}}, \frac{1}{1}) = 1$

Basic var.	$x_1$	$x_2$	$x_3$	$x_4$	RHS
$z$	0	0	1	0	4
$x_2$	0	1	$\frac{1}{2}$	$-\frac{1}{2}$	1.5
$x_1$	1	0	0	1	1

$B = \{1, 2\}, N = \{3, 4\}$

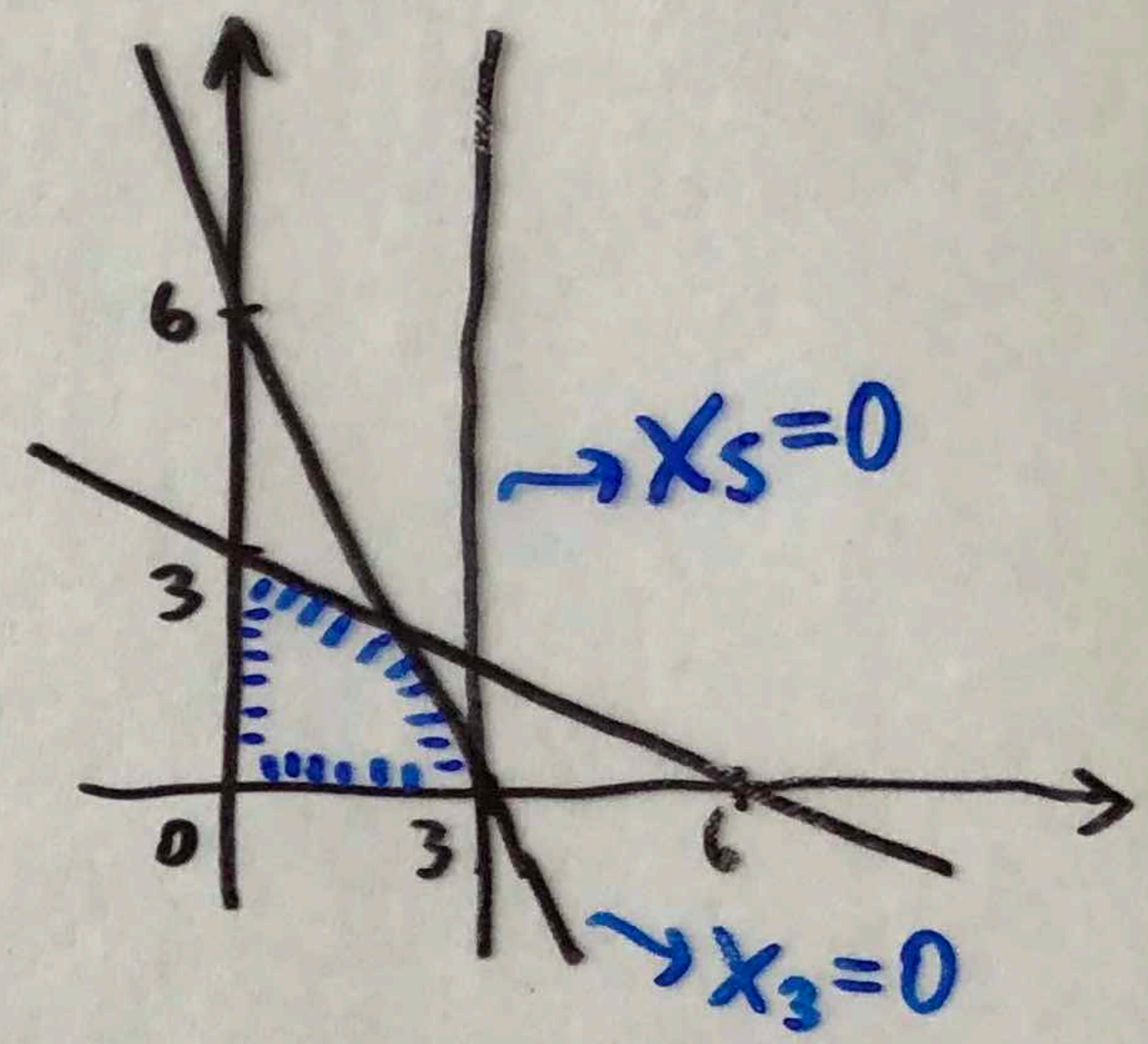
$\tilde{x}^* = \begin{pmatrix} 1 \\ 1.5 \\ 0 \\ 0 \end{pmatrix}$  is another optimal sol-n

In fact, any point on the segment joining  $\begin{pmatrix} 0 \\ 2 \\ 0 \\ 1 \end{pmatrix}$  and  $\begin{pmatrix} 1 \\ 1.5 \\ 0 \\ 0 \end{pmatrix}$  is optimal.

# Breaking Ties, Degeneracy

- Tie for entering variable  $\rightarrow$  break the tie arbitrarily
- Tie in the min-ratio test

Ex.:  $\max X_1 + X_2$   
 s.t.  $2X_1 + X_2 \leq 6$   
 $2X_1 + 4X_2 \leq 12$   
 $X_1 \leq 3$   
 $X_1, X_2 \geq 0$



The augmented system:

$$\begin{aligned} 2X_1 + X_2 + X_3 &= 6 \\ 2X_1 + 4X_2 + X_4 &= 12 \\ X_1 + X_5 &= 3 \\ X_i &\geq 0 \quad \forall i \end{aligned}$$

$X_2, X_5$  nonbasic  $\Rightarrow \begin{pmatrix} X_1 \\ X_3 \\ X_4 \end{pmatrix} = \begin{pmatrix} 3 \\ 0 \\ 6 \end{pmatrix}$   
 $X_2, X_3$  nonbasic  $\Rightarrow \begin{pmatrix} X_1 \\ X_4 \\ X_5 \end{pmatrix} = \begin{pmatrix} 3 \\ 6 \\ 0 \end{pmatrix}$   
 $X_3, X_5$  nonbasic  $\Rightarrow \begin{pmatrix} X_1 \\ X_2 \\ X_4 \end{pmatrix} = \begin{pmatrix} 3 \\ 0 \\ 6 \end{pmatrix}$

all three basic solutions:

$$x = \begin{pmatrix} 3 \\ 0 \\ 0 \\ 6 \\ 0 \end{pmatrix}$$

Def-n: A basic sol-n  $X$  is called degenerate if it has a basic variable that is 0. (5)

Continue with the example.

Basic var.	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	RHS
$Z$	-1	-1	0	0	0	0
$x_3$	2*	1	1	0	0	6
$x_4$	2	4	0	1	0	12
$x_5$	1	0	0	0	1	3

min-ratio test:  $\min\left(\frac{6}{2}, \frac{12}{2}, \frac{3}{1}\right) = \min(3, 6, 3) = 3$

↑ tie ↑

2 basic variables  $x_3$  and  $x_5$  drop to 0 simultaneously; but only one of them can become nonbasic.

Thus, tie in min-ratio test  $\rightarrow$  next basic sol-n is degenerate:

Basic var.	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	RHS
$Z$	0	-0.5	0.5	0	0	3
$x_1$	1	0.5	0.5	0	0	3
$x_4$	0	3	-1	1	0	6
$x_5$	0	-0.5	-0.5	0	1	0

- (6)
- What problems might degeneracy cause?

In degenerate LP's, Simplex method might cycle: repeat the same sequence of solutions without increasing the obj. f-n. (see Beale's example)

- To prevent cycling, "anti-cycling rules" are invented.

- One of the best-known rules is by Prof. Robert Bland (Cornell OR&IE)

Least-index rule ("Bland's rule"):

i) Among those variables eligible to enter the basis, choose the one with the smallest index (the one which is most to the left in the tableau).

ii) Among all variables eligible to leave the basis, select the one with the smallest index.

E.g.: Suppose  $x_6$ ,  $x_2$  and  $x_4$  are eligible to leave (ratios for corresponding rows are the same)  $\rightarrow$  choose  $x_2$ .