Math 4620/5620

Homework 1 Solution

1. There are many different ways of doing these formulations, and here we present just one of them.

Decision Variables:

Let x_{ci} , y_{ci} and z_{ci} represent the number of flights per day made to city c by the *i*'th plane of B737, AB321 and ATR-70 respectively.

Constraints:

1. minimum trip constraints

$$\sum_{i=1}^{10} x_{\text{LA},i} + \sum_{i=1}^{15} y_{\text{LA},i} = 2 \text{ trips}$$

$$\sum_{i=1}^{10} x_{\text{BOS},i} + \sum_{i=1}^{15} y_{\text{BOS},i} + \sum_{i=1}^{2} z_{\text{BOS},i} \ge 4 \text{ trips}$$

$$\sum_{i=1}^{10} x_{\text{DC},i} + \sum_{i=1}^{15} y_{\text{DC},i} + \sum_{i=1}^{2} z_{\text{DC},i} \ge 4 \text{ trips}$$

$$\sum_{i=1}^{10} x_{\text{MIA},i} + \sum_{i=1}^{15} y_{\text{MIA},i} \ge 4 \text{ trips}$$

And of course, non-negativity is important as well.

 $x_{ci}, y_{ci}, z_{ci} \ge 0$ for all city c and plane i

2. maximum flying time constraints

$$\begin{array}{rcl} x_{BOS,i} + 2x_{DC,i} + 4x_{MIA,i} + 6x_{LA,i} &\leq 15 \text{ hours for each } i = 1, \cdots, 10 \\ y_{BOS,i} + 2y_{DC,i} + 5y_{MIA,i} + 8y_{LA,i} &\leq 15 \text{ hours for each } i = 1, \cdots, 15 \\ 2z_{BOS,i} + 4z_{DC,i} &\leq 18 \text{ hours for each } i = 1, 2 \end{array}$$

Objective Functions:

1.

$$\begin{array}{ll}\text{Minimize} & 6000 \sum_{i=1}^{10} x_{BOS,i} + 7000 \sum_{i=1}^{10} x_{DC,i} + 8000 \sum_{i=1}^{10} x_{MIA,i} + 10000 \sum_{i=1}^{10} x_{LA,i} + \\ & 4000 \sum_{i=1}^{15} y_{BOS,i} + 3500 \sum_{i=1}^{15} y_{DC,i} + 6000 \sum_{i=1}^{15} y_{MIA,i} + 10000 \sum_{i=1}^{15} y_{LA,i} + \\ & 1000 \sum_{i=1}^{2} z_{BOS,i} + 2000 \sum_{i=1}^{2} z_{DC,i} \end{array}$$

$$\begin{aligned} \text{Maximize} \qquad & -1000\sum_{i=1}^{10}x_{BOS,i} + 2000\sum_{i=1}^{10}x_{MIA,i} + 8000\sum_{i=1}^{10}x_{LA,i} + \\ & -1000\sum_{i=1}^{15}y_{BOS,i} + 2000\sum_{i=1}^{15}y_{DC,i} + 2000\sum_{i=1}^{15}y_{MIA,i} + 4000\sum_{i=1}^{15}y_{LA,i} + \\ & 2000\sum_{i=1}^{2}z_{BOS,i} + 2000\sum_{i=1}^{2}z_{DC,i} \end{aligned}$$

3.

$$\begin{aligned} \text{Minimize} \qquad & 1\sum_{i=1}^{10} x_{BOS,i} + 2\sum_{i=1}^{10} x_{DC,i} + 4\sum_{i=1}^{10} x_{MIA,i} + 6\sum_{i=1}^{10} x_{LA,i} + \\ & 1\sum_{i=1}^{15} y_{BOS,i} + 2\sum_{i=1}^{15} y_{DC,i} + 5\sum_{i=1}^{15} y_{MIA,i} + 8\sum_{i=1}^{15} y_{LA,i} + \\ & 2\sum_{i=1}^{2} z_{BOS,i} + 4\sum_{i=1}^{2} z_{DC,i} \end{aligned}$$

2. Let x_{pi} where p = A, B, C and i = 1, 2, 3, 4 be the amount (MG / day) of water treated in plant p using treatment i. Then we want to minimize the total cost

$$\sum_{i=1}^{4} c_i (x_{Ai} + x_{Bi} + x_{Ci})$$

subject to the following set of constraints:

lower bound:

$$\sum_{i=1}^{4} x_{pi} \ge K_p \text{ for each } p = A, B, C$$

upper bound (we cannot treat more water than the flow):

$$\sum_{i=1}^{4} x_{pi} \le M_p \text{ for each } p = A, B, C$$

where $M_A = M_1$, $M_B = M_2$ and $M_C = M_1 + M_2$

2.

permissible pollutant:

$$\sum_{i=1}^{4} x_{Ai} \cdot P_i \leq P_S \cdot M_1 \text{ (Phenol at A)}$$

$$\sum_{i=1}^{4} x_{Ai} \cdot N_i \leq N_S \cdot M_1 \text{ (Nitrogen at A)}$$

$$\sum_{i=1}^{4} x_{Bi} \cdot P_i \leq P_B \cdot M_2 \text{ (Phenol at B)}$$

$$\sum_{i=1}^{4} x_{Bi} \cdot N_i \leq N_B \cdot M_2 \text{ (Nitrogen at B)}$$

$$\sum_{i=1}^{4} (x_{Ai} + x_{Bi} + x_{Ci}) \cdot P_i \leq P_S \cdot (M_1 + M_2) \text{ (Phenol at C)}$$

$$\sum_{i=1}^{4} (x_{Ai} + x_{Bi} + x_{Ci}) \cdot N_i \leq N_S \cdot (M_1 + M_2) \text{ (Nitrogen at C)}$$

nonnegativity:

$$x_{pi} \ge 0$$
 for each $p = A, B, C$, and $i = 1, 2, 3, 4$

3. Let x, y and z denote the number of kilograms of corn, tankage and alfalfa respectively. Then we have the following LP:

4. (a) TRUE. Consider the isoprofit line through (0, 2), the isoprofit line through (3, 3), and the isoprofit line through (6, 3). Recall that these lines are all parallel to each other. Since (3, 3) produces a larger value of the objective function than (0, 2) and (6, 3), it follows that the isoprofit line through (3, 3) cannot coincide with the isoprofit line through (0, 2) or (6, 3), and moreover, it can be obtained by shifting (while maintaining the slope) the isoprofit line through (0, 2) or (6, 3) in the direction of increasing z. This means that (3, 3) corresponds to the isoprofit line with maximum z.

An example of objective function: $Z = 6x_2 - x_1$.

(b) TRUE. Since (3,3) is an optimal solution, it is clear from the graph that multiple optimal solutions will only exist if the isoprofit line through (3,3) coincides with either one of the boundaries (of the feasible region) containing (3,3). Hence, either (0,2) or (6,3) are also optimal. Examples of objective functions: Z = 3x₂ - x₁ (makes (0,2) optimal), Z = x₂ (makes (6,3))

optimal).

(c) FALSE. Consider the objective function $z = -x_1 - x_2$.

5. (a) The LP is:

 $\begin{array}{rll} \max & 16x_1 + 20x_2 & (\text{revenue}) \\ \text{s. t.} & 3x_1 + 4x_2 & \leq & 12 \text{ (skilled labor)} \\ & 4x_1 + 2x_2 & \leq & 8 \text{ (unskilled labor)} \\ & 2x_1 + x_2 & \leq & 16 \text{ (raw material)} \\ & x_2 & \geq & 2 \text{ (marketing consideration)} \\ & x_1, x_2 & \geq & 0 \end{array}$

(b) The feasible region is the quadrangle with the vertices at (0,2), (0,3), (1,2) and (0.8,2.4). The optimal solution is (0.8,2.4) with optimal value 60.8.

6.

Decision variables.

Let

 x_t be the number of calculators produced in time period t during regular shift $t = 1, \dots, 6$

 y_t be the number of calculators produced in time period t during overtime hours $t=1,\ldots,6$

 i_t be the number of calculators stored at the end of period t

t = 0,...,6

Objective function:

$$Min \sum_{t=1}^{6} (10x_t + 12y_t + 0.5i_t)$$

Subject to

Production capacity constraints:

 $\begin{array}{ll} x_t \leq 100 & \mbox{ for } t = 1, \dots, 6 \\ y_t \leq 75 & \mbox{ for } t = 1, \dots, 6 \end{array}$

No inventory at the beginning of time period 1:

 $i_0 = 0$

Balance constraints:

 $i_{t-1} + x_t + y_t = d_t + i_t$ for t = 1,...,6

(i.e.,

inventory from previous time period + total number produced in this time period = demand in current time period + inventory at the end of current time period)

Nonnegativity constraints:

 $x_t, y_t, i_t \ge 0 \qquad \qquad \text{for } t = 1, \dots, 6$